

## THE EXPRESSION PROBLEM

**Exercise 1.** Implement the `string_of_` functions or methods, covering all data cases, corresponding to the `eval_` functions in at least two examples from the lecture, including both an object-based example and a variant-based example (either standard, or polymorphic, or extensible variants).

**Exercise 2.** Split at least one of the examples from the previous exercise into multiple files and demonstrate separate compilation.

**Exercise 3.** Can we drop the tags `Lambda_t`, `Expr_t` and `LExpr_t` used in the examples based on standard variants (file `FP_ADT.ml`)? When using polymorphic variants, such tags are not needed.

**Exercise 4.** Factor-out the sub-language consisting only of variables, thus eliminating the duplication of tags `VarL`, `VarE` in the examples based on standard variants (file `FP_ADT.ml`).

**Exercise 5.** Come up with a scenario where the extensible variant types-based solution leads to a non-obvious or hard to locate bug.

**Exercise 6.** \* Re-implement the direct object-based solution to the expression problem (file `Objects.ml`) to make it more satisfying. For example, eliminate the need for some of the `rename`, `apply`, `compute` methods.

**Exercise 7.** Re-implement the visitor pattern-based solution to the expression problem (file `Visitor.ml`) in a functional way, i.e. replace the mutable fields `subst` and `beta_redex` in the `eval_lambda` class with a different solution to the problem of treating `abs` and non-`abs` expressions differently.

\* See if you can replace the reference cells `result` in `evalN` and `freevarsN` functions (for  $N = 1, 2, 3$ ) with a different solution to the problem of polymorphism wrt. the type of the computed values.

**Exercise 8.** Extend the sub-language `expr_visit` with variables, and add to arguments of the evaluation constructor `eval_expr` the substitution. Handle the problem of potentially duplicate fields `subst`. (One approach might be to use ideas from exercise 6.)

**Exercise 9.** Implement the following modifications to the example from the file `PolyV.ml`:

1. Factor-out the sub-language of variables, around the already present `var` type.
2. Open the types of functions `eval3`, `freevars3` and other functions as required, so that explicit subtyping, e.g. in `eval3 [] (test2 :=> lexpr_t)`, is not necessary.
3. Remove the double-dispatch currently in `eval_lexpr` and `freevars_lexpr`, by implementing a cascading design rather than a “divide-and-conquer” design.

**Exercise 10.** Streamline the solution `PolyRecM.ml` by extending the language of  $\lambda$ -expressions with arithmetic expressions, rather than defining the sub-languages separately and then merging them. See slide on page 15 of Jacques Garrigue *Structural Types, Recursive Modules, and the Expression Problem*.

**Exercise 11.** Transform a parser monad, or rewrite the parser monad transformer, by adding state for the line and column numbers.

\* How to implement a monad transformer transformer in OCaml?

**Exercise 12.** Implement `_of_string` functions as parser combinators on top of the example `PolyRecM.ml`. Sections 4.3 and 6.2 of *Monadic Parser Combinators* by Graham Hutton and Erik Meijer might be helpful. Split the result into multiple files as in Exercise 2 and demonstrate dynamic loading of code.

**Exercise 13.** What are the benefits and drawbacks of our lazy-monad-plus (built on top of *odd lazy lists*) approach, as compared to regular monad-plus built on top of *even lazy lists*? To additionally illustrate your answer:

1. Rewrite the parser combinators example to use regular monad-plus and even lazy lists.
2. Select one example from Lecture 8 and rewrite it using lazy-monad-plus and odd lazy lists.