

A BOUND ON SOLUTIONS OF LINEAR INTEGER EQUALITIES AND INEQUALITIES

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ABSTRACT. Consider a system of linear equalities and inequalities with integer coefficients. We describe the set of rational solutions by a finite generating set of solution vectors. The entries of these vectors can be bounded by the absolute value of a certain subdeterminant. The smallest integer solution of the system has coefficients not larger than this subdeterminant times the number of indeterminates. Up to the latter factor, the bound is sharp.

Let A, B, C, D be $m \times n$ -, $m \times 1$ -, $p \times n$ -, $p \times 1$ -matrices respectively with integer entries. The rank of A is r , and s is the rank of the $(m+p) \times n$ -matrix $\begin{pmatrix} A \\ C \end{pmatrix}$. Let M be an upper bound on the absolute values of those $(s-1) \times (s-1)$ - or $s \times s$ -subdeterminants of the $(m+p) \times (n+1)$ -matrix $\begin{pmatrix} A & B \\ C & D \end{pmatrix}$, which are formed with at least r rows from (A, B) .

THEOREM. *If $Ax = B$ and $Cx \geq D$ have a common integer solution, then they have one with coefficients bounded by $(n+1)M$.*

Let M_1, M_2 , and M_3 be upper bounds on the absolute values of the $r \times r$ -subdeterminants, the subdeterminants, and the entries of (A, B) respectively. Taking the $n \times n$ -identity matrix for C and $D = 0$, we have the following

COROLLARY. *If $Ax = B$ has a nonnegative integer solution, then it has one with coefficients bounded by $(n+1)M_1$.*

S. Cook [4] obtained a bound of the order of $M_3^{m^2}$ in this case. I. Borosh and L. B. Treybig conjecture that one can always have the bound M_1 . For many cases, this bound would be sharp. They give an elegant proof for M_2^2 in [2]. In [1], [3] they obtain M_1 in the cases where $r = n-1$ and for homogeneous systems (only nontrivial solutions being considered), and nrM_1 if the matrix has no $r \times r$ -subdeterminants which are zero. Their work arose from topological questions, while Cook's and our aim was to prove that the solvable linear integer programs form a NP -complete set (see Remarks 2 and 3).

For the proof of the theorem we first note that it suffices to consider the case $s = n$. For if $s < n$, then choose an integer solution y , let e_i be 1 or -1 according to whether y_i is ≥ 0 or < 0 . To the given system add $n-s$

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inequalities $e_i x_i \geq 0$ so that the resulting system has rank n . Then the theorem for this system implies that for the one with rank s .

Let $S = [A = B] \cap [C \geq D] = \{x \in \mathbb{Q}^n: Ax = B \text{ and } Cx \geq D\}$, $H = [A = 0] \cap [C \geq 0]$. The i th row A_i of A corresponds to a linear form on \mathbb{Q}^n , which we also denote by A_i , and similarly for C . Let $I \subseteq \{1, \dots, p\}$. If $S \cap \bigcap_{i \in I} [C_i = D_i]$ is nonempty and the matrix with rows A_1, \dots, A_m and C_i for $i \in I$ has rank n , then the single vector in this set is called a vertex of S . We call $H \cap \bigcap_{i \in I} [C_i = 0]$ an edge of H , if it is not equal to $\{0\}$, and the matrix with rows A_1, \dots, A_m and C_i for $i \in I$ has rank $n - 1$. We may assume $\#I = n - r$ for a vertex, and $\#I = n - r - 1$ for an edge. Choose a nonzero vector in each edge of H . Let F be the set of all these vectors, and E the set of all vertices of S .

LEMMA. $S = \text{conv}(E) + \mathbb{Q}_+ F$.

Here $\text{conv}(E) + \mathbb{Q}_+ F$ denotes the set of all $\sum_{e \in E} \lambda_e e + \sum_{f \in F} \mu_f f$ with nonnegative rational numbers λ_e , μ_f , and $\sum_{e \in E} \lambda_e = 1$. The proof is by induction on n . $n = 0$ is trivial, the empty sum being 0. Obviously $\text{conv}(E) + \mathbb{Q}_+ F \subset S$. So let $n > 0$, $x \in S$.

Case 1. $r > 0$.

Suppose that $A_{1k} \neq 0$, and let $u: [A_1 = 0] \rightarrow \mathbb{Q}^{n-1}$ be the projection that takes the k th coordinate to zero. u is an isomorphism of vector spaces. $A'_j = A_j \circ u^{-1}$ and $C'_i = C_i \circ u^{-1}$ are linear forms on \mathbb{Q}^{n-1} , and the rank of the matrix $\begin{pmatrix} A'_j \\ C'_i \end{pmatrix}$ is $n - 1$. Let $B'_j = 0$, $D'_i = D_i - C_i(x)$, $E' = u(E - x)$, $F' = u(F)$. Here $E - x = \{e - x: e \in E\}$. We show that E' contains every vertex and F' a nonzero vector of each edge with respect to A' , B' , C' , D' . Let $\{y\} = [A' = B'] \cap \bigcap_{i \in I} [C'_i = D'_i]$ be a vertex of $[A' = B'] \cap [C' \geq D']$. Then $\{u^{-1}(y) + x\} = [A = B] \cap \bigcap_{i \in I} [C_i = D_i]$ is a vertex of S , and hence $u^{-1}(y) + x \in E$. Then $y \in E'$. Let $Z = [A' = 0] \cap [C' \geq 0] \cap \bigcap_{i \in I} [C'_i = 0]$ be an edge of $[A' = 0] \cap [C' \geq 0]$. Then $u^{-1}(Z) = H \cap \bigcap_{i \in I} [C_i = 0]$ is an edge of H , and hence there is an $f \in u^{-1}(Z) \cap F \setminus \{0\}$. Then $u(f) \in Z \cap F' \setminus \{0\}$.

We have $0 \in [A' = B'] \cap [C' \geq D']$, and using the induction hypothesis, we obtain $0 \in \text{conv}(E') + \mathbb{Q}_+ F'$,

$$x \in \text{conv}(E) + \mathbb{Q}_+ F.$$

Case 2. $r = 0$.

We assume $B = 0$ and $n > 1$. Then for some i, j , C_i and C_j are linearly independent. We can find a $y_1 \in [C_j > 0] \cap [C_i < 0]$. Let $\lambda > D_i/C_i(y_1)$, $(2C_j(x) - D_j)/C_j(y_1)$. Then for $y = \lambda y_1$ we have $C_i(y) < D_i$ and $C_j(2x - y) < D_j$, and hence the intersection of $[C \geq D]$ with the line

$$l = \{x + \alpha(y - x): \alpha \in \mathbb{Q}\}$$

is a finite interval $[z_1, z_2]$ containing x . Thus it is sufficient to prove that $z_1, z_2 \in \text{conv}(E) + \mathbb{Q}_+ F$. Let z be z_1 or z_2 . There is a j with $z \in [C_j = D_j]$, for

otherwise an open interval on l containing z would be contained in $[C \geq D]$. Suppose that $C_{jk} \neq 0$, and let $u: [C_j = 0] \rightarrow \mathbb{Q}^{n-1}$ be the projection that takes the k th coordinate to zero. Let $C'_i = C_i \circ u^{-1}$ and $D'_i = D_i - C_i(z)$, $E' = u((E - z) \cap [C_j = 0])$, $F' = u(F \cap [C_j = 0])$. To a vertex of $[C' \geq D']$ with index set I corresponds the vertex of $[C \geq D]$ with index set $I \cup \{j\}$, and similarly for edges. Like in Case 1, it now follows that $z \in \text{conv}(E) + \mathbb{Q}_+ F$. This completes the proof of the lemma.

We can now prove the theorem stated at the beginning. By Cramer's rule, we can choose the sets E and F such that all coefficients of their elements are bounded by M , and all elements of F are integer vectors. Let x be an integer solution of $Ax = B$ and $Cx \geq D$. Then $x \in \text{conv}(E) + \mathbb{Q}_+ F$, and it is easy to see [6, 2.2] that there exists $F' \subseteq F$ with $\#F' \leq n$ and $x \in \text{conv}(E) + \mathbb{Q}_+ F'$. Take a representation

$$x = \sum_{e \in E} \lambda_e e + \sum_{f \in F'} \mu_f f.$$

Then $x - \sum_{f \in F'} [\mu_f] f$ is the desired integer solution of $Ax = B$ and $Cx \geq D$ with coefficients bounded by $\sum_E \lambda_e M + \sum_{F'} M \leq (n+1)M$.

REMARKS. 1. The lemma states a decomposition of the set S of rational solutions into the sum of a compact convex set and a pointed cone (i.e. one that does not contain a line). A corresponding decomposition can be made without the assumption that $s = n$, by adding the common null space of A and C . However, the construction of the generating set F for the cone requires a little more care, and the bounds obtained in this way are considerably worse than the one in the theorem.

2. It follows from the theorem that the set of all A, B, C, D as above, for which there exists an integer vector x with $Ax = B$ and $Cx \geq D$, is in NP . It is then easy to see that it is an NP -complete set. The same holds for slight variations of the problem, such as $\{(A, B): Ax = B \text{ has a nonnegative integer solution}\}$, $\{(C, D): Cx \geq D \text{ has an integer solution}\}$, and $\{(C, D): Cx \geq D \text{ has a nonnegative integer solution}\}$. One easily finds polynomial transformations between these sets. For terminology concerning NP , see [5], e.g.

3. For any entry C_{ij} of C , let $C_{ij} = \sum_k C'_{ijk} \cdot 2^k$ with $C'_{ijk} \in \{-1, 0, 1\}$, and for $x \in \mathbb{Q}^n$ let $y_{jk} = x_j \cdot 2^k$. Then

$$Cx \geq D \Leftrightarrow \text{for all } i, \sum_{j,k} C'_{ijk} y_{jk} \geq D_i.$$

Furthermore, for any j

$$\text{for all } k, y_{jk} = x_j \cdot 2^k \Leftrightarrow$$

$$y_{j0} = x_j \text{ and for all } k \text{ there exists } z_{jk} \in \mathbb{Q}^n \text{ with}$$

$$z_{jk} = y_{j,k-1} \text{ and } y_{jk} = y_{j,k-1} + z_{jk}.$$

Combining these, we obtain a polynomial transformation from the set of all solvable systems of linear integer inequalities to the set of those inequalities

(C, D) , for which all entries of C are $-1, 0$, or 1 . This set is therefore NP -complete. For C consisting of 0 's and 1 's only, the problem is in P .

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