



Interactive Proof Tools Assignment 6

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<http://www.mpi-sb.mpg.de/~nivelle/teaching/intprooftools2003/main.html>

The exercises 6.1 to 6.3 belong to our project of formalizing a (propositional) resolution calculus in PVS.

Exercise 6.1 A (propositional) *clause* is a finite disjunction of literals, where a *literal* is either a propositional variable or the negation of a propositional variable.

1. Formalize clauses in PVS, i.e., define non-empty types of propositional variables, literals and clauses.

Hint: A clause can be represented as a multiset of literals.

2. Define a function `occurs` which maps a set of clauses to the set of its propositional variables.

Hint: First define a function which maps a single clause to the set of its variables.

3. Prove that `occurs` maps a finite set of clauses to a finite set of variables.

Exercise 6.2 An *interpretation* is a mapping of propositional variables to *truth values*.

1. Define a type of truth values and a type of interpretations.
2. Define a predicate `true?` which holds of a clause and an interpretation iff the clause evaluates to true in the interpretation.
3. Define predicates `valid?` and `satisfiable?` which hold of a clause if it is *valid* (true in all interpretations) or *satisfiable* (true in at least one interpretation), respectively.
4. Prove that the clauses $A \vee \neg A$ resp. $A \vee \neg B$ (with propositional variables A, B) are valid resp. satisfiable.
5. Extend the definitions of validity and satisfiability to sets of clauses. Prove that validity of a set of clauses implies satisfiability.

Exercise 6.3 Prove the following theorem: Given a clause c and two interpretations I_1 and I_2 , if I_1 and I_2 agree on all propositional variables occurring in c then c is true in I_1 iff c is true in I_2 .

Exercise 6.4 Prove the following formulas in Coq:

1. $(x, y, z: \text{nat}) (1e\ x\ y) \wedge (1e\ y\ z) \rightarrow (1e\ x\ z)$
2. $(x, y: \text{nat}) (1e\ x\ y) \vee (1e\ y\ x)$